

# About the Abel map computation

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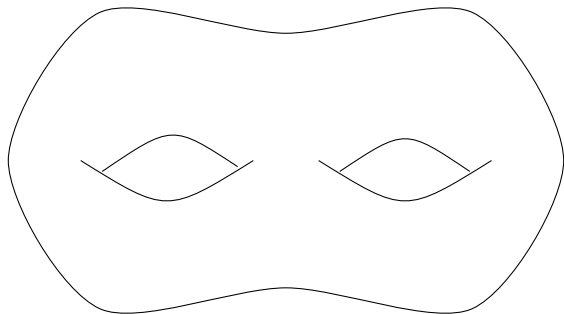
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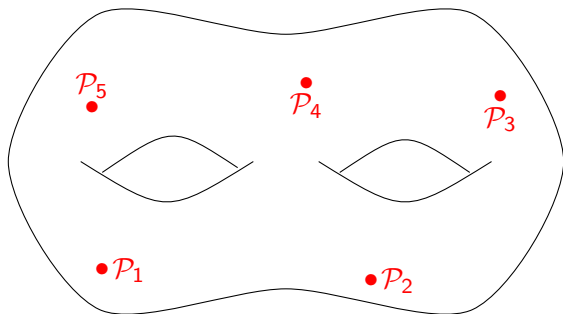
## Divisor of a Riemann surface

- $\mathbb{K} = \mathbb{Q}(\alpha)$  ;  $F \in \mathbb{K}[X, Y]$  defines a Riemann surface of genus  $g$ .



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- Divisor  $D = \sum n_{\mathcal{P}} \mathcal{P}$  formal sum of points.
- Degree of a divisor:  $\deg(D) = \sum n_{\mathcal{P}}$
- Function divisor  $(f) = \sum v_{\mathcal{P}}(f) \mathcal{P} \rightarrow$  zeros and poles of  $f$ .

## Is a degree 0 divisor a function divisor ?

- Let
- $\mathcal{F}$  the set of function divisors.
  - $\mathcal{D}_0$  the set of degree 0 divisors.

We have  $\mathcal{F} \subset \mathcal{D}_0$  but no equality. Thus the following questions:

- Let  $D \in \mathcal{D}_0$ . Does  $D \in \mathcal{F}$  ?
- Is there  $n \in \mathbb{Z}$  such that  $nD \in \mathcal{F}$  ?
- Given  $D_1, \dots, D_p \in \mathcal{D}_0$ , is there  $n_1, \dots, n_p \in \mathbb{Z}$   
s.t.  $n_1 D_1 + \dots + n_p D_p \in \mathcal{F}$  ?

## One answer: Abel-Jacobi theorem

### Theorem

*The Abel map*

$$\begin{aligned} \mathcal{A}: \mathcal{D}_0/\mathcal{F} &\longrightarrow \mathbb{C}^g/\Gamma \\ \sum n_{\mathcal{P}}\mathcal{P} &\longmapsto \sum n_{\mathcal{P}} \left( \int_{\mathcal{O}}^{\mathcal{P}} \omega_1, \dots, \int_{\mathcal{O}}^{\mathcal{P}} \omega_g \right) \end{aligned}$$

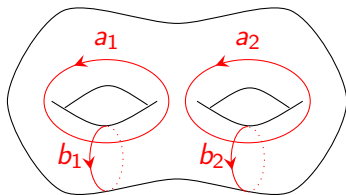
*is a group isomorphism.*

$\implies$  Computing the Abel map answers our questions !

## Details on the Abel map

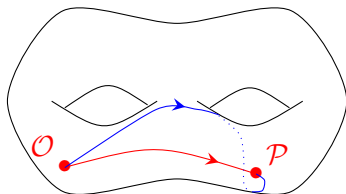
$$\begin{aligned} \mathcal{A}: \mathcal{D}_0/\mathcal{F} &\longrightarrow \mathbb{C}^g/\Gamma \\ \sum n_P \mathcal{P} &\longmapsto \sum n_P \left( \int_O^{\mathcal{P}} \omega_1, \dots, \int_O^{\mathcal{P}} \omega_g \right) \end{aligned}$$

- There are  $2g$  cycles that cannot be reduced to a point.  
 $\implies$  homology basis.



- Periods: integrals of holomorphic differentials over this basis.

- Modulo the periods,  $\int_O^{\mathcal{P}} \omega_i$  does not depend on the path choice.



## Computer algebra

- Integration (logarithmic part)

Risch 69, *The problem of integration in finite terms*

Bronstein 1990, *Integration of elementary functions*

Bertrand 95, *Computing a hyperelliptic integral using arithmetic in the jacobian of the curve*

- Symbolic ODEs

- algebraic solutions,

Baldassari & Dwork 79, *On second order linear differential equations with algebraic solutions*

- Galois group

Compoint & Singer 79, *Relations linéaires entre solutions d'équations différentielles*

## Physics

- Quasi periodic solutions of the KP equation
- Special solutions of the KdV and NLS equations

Deconinck & Segur 1998, *The KP equation with quasiperiodic initial data*

Deconinck & Patterson 2007, *Computing the Abel map*

# Computing the Abel map: overview

Patterson 2007, *Algebra-geometric algorithms for integrable systems*

## ① Computing the period matrix

van Hoeij & Deconinck 2001, *Computing Riemann matrices of algebraic curves*

- Monodromy computation,
- Use it to compute the homology basis,
  - ⇒ Tretkoff<sup>2</sup> 1984, *Combinatorial Group Theory, Riemann Surfaces and Differential Equations*.
- Compute a basis of holomorphic differential on the surface,
- Compute the periods.

## ② Make analytic continuation along the needed paths.



## Regular and critical points: a projective point of view

$$\mathcal{C} = \{(x, y) \in \mathbb{C}^2 \mid F(x, y) = 0\}$$

Let  $x_0 \in \mathbb{C}$ :

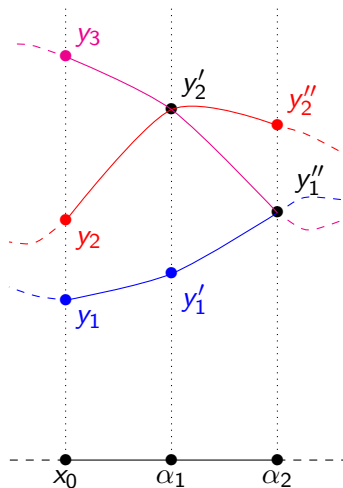
- **Fiber** at  $x_0$ :

$$\mathcal{F}(x_0) = \{\text{roots of } F(x_0, Y) = 0\}.$$

- **Regular point** :  $\#\mathcal{F}(x_0) = d_Y$ .

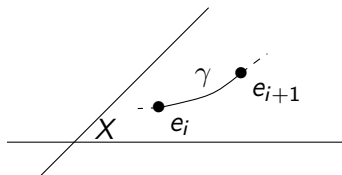
- **Critical point** :  $\#\mathcal{F}(x_0) < d_Y$ .

$$\implies \text{roots of } R_F = \text{Res}_Y(F, F_Y)$$

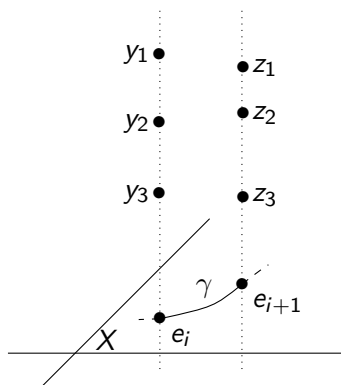


## Analytic continuation : compute fibers and connect

- 1 Choice of paths,
- 2 Choice of connection points,

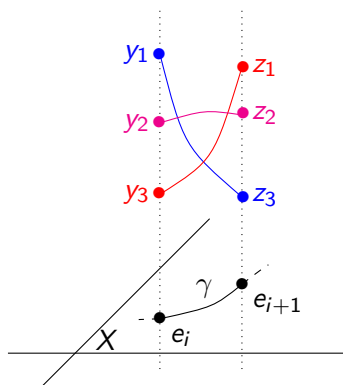


## Analytic continuation : compute fibers and connect



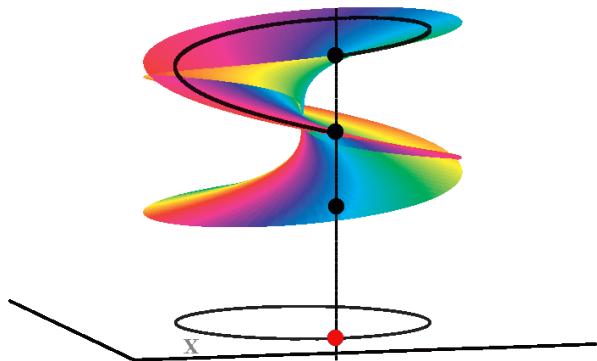
- 1 Choice of paths,
- 2 Choice of connection points,
- 3 Compute the fibers,

## Analytic continuation : compute fibers and connect



- 1 Choice of paths,
- 2 Choice of connection points,
- 3 Compute the fibers,
- 4 Connection method.

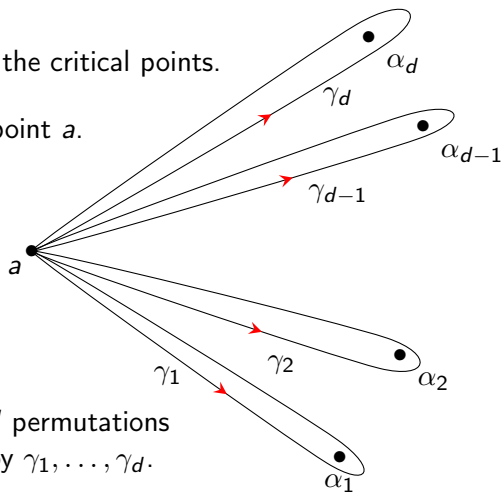
Local monodromy:  $F(X, Y) = Y^3 - X$  around  $X = 0$



- analytic continuation around the critical point  $X = 0$   
 $\implies$  permutation on the fiber

# Monodromy group

- Let  $\alpha_1, \dots, \alpha_d$  denote the critical points.
- We fix a regular base point  $a$ .



- We will compute the  $d$  permutations  $\sigma_1, \dots, \sigma_d$  generated by  $\gamma_1, \dots, \gamma_d$ .
- These permutations generate the monodromy group.

# Computing the Abel map: details

Patterson 2007, *Algebro-geometric algorithms for integrable systems*

## 1 Computing the period matrix

van Hoeij & Deconinck 2001, *Computing Riemann matrices of algebraic curves*

- Monodromy computation,
  - Use it to compute the homology basis,
  - Compute a basis of holomorphic differential on the surface,
  - Compute the periods.
- ## 2 Make analytic continuation along the needed paths.
- Makes analytic continuation using first derivative only,
  - Uses Puiseux expansion when  $\mathcal{P}$  is close from a critical point.

# Connexion method

- van Hoeij & Deconinck 2001
  - Connexion using only derivatives (order 1),
  - No certification that the connexion is correct.
- van Hoeij & Rybowicz 2006
  - Certified algorithm using order 1 truncations,
  - Efficiency problem: sometimes too many intermediary steps (two close critical points).
- Poteaux 2007, *Computing monodromy groups defined by plane algebraic curves*
  - Power series expansions at well chosen truncation orders,
  - Trade-off between truncations orders and the number of steps,
  - (similar to Chudnovsky<sup>2</sup> 86, 87, 90, van der Hoeven 99 for ODEs)
  - *useful for the integration part.*



## Choice of the paths

- Most methods: going as far as possible from the critical points (one loop = 17 steps)  
van Hoeij & Deconinck 01, Chudnovsky<sup>2</sup> 86, 87, 90, van der Hoeven 99
  - numerical problems when two critical points are very close. . .
- Another way : computing series *above critical points*
  - van der Hoeven 01 (ODE) Poteaux 07 (Puiseux series)
  - one loop = one evaluation,
  - Puiseux series are also useful for the Abel map !
    - needed when the point  $\mathcal{P}$  is close to a critical point,  
Deconinck-Patterson 2007, *Computing the Abel map*
    - it enables to compute integral basis (thus cohomology).  
van Hoeij 1994, *an algorithm for computing an integral basis in an algebraic function field*

## Puiseux series

- A generalization of power series:

- above a regular point:  $Y(X) = \sum_{k=0}^{\infty} \beta_k (X - x_0)^k$

- above a critical point:  $Y(X) = \sum_{k=n}^{\infty} \beta_k (X - x_0)^{k/e}$

- Costly to compute:

- coefficient growth and high degree extension field
- one example:  $(Y^3 - X)((Y - 1)^2 - X)(Y - 2 - X^2) + X^2 Y^5$   
 $\implies$  extension of degree 92, rational numbers with 132 digits
- Bit complexity:  $\mathcal{O}(d_Y^{32} d_X^4 \text{ht}(F)^2)$  Walsh 2000
- Best known arithmetic complexity:  $\mathcal{O}(D^5)$  Poteaux 2008

# Computing Puiseux series efficiently

- A **modular-numeric strategy** to avoid coefficient growth:
  - First computation modulo a *well-chosen* prime number  
 $\implies$  structure of the solution,
  - Using it, conduct numerical computations to get an approximation of the coefficients.
- Towards a **faster algorithm**
  - Computing fast “Taylor parts” + factorize  $F$  in  $\mathbb{K}[[X]][Y]$  during the process + using optimal truncation bounds via relaxed algorithms  $\implies \mathcal{O}(D^4)$
  - Being able to factorize  $F = F_1 \cdot F_2$  in  $\mathbb{K}[[X]][Y]$  knowing singular part of Puiseux series of  $F_1 \implies \mathcal{O}(D^3)$

## What to be done yet ?

- Numerical part of Puiseux is not certified and might be improved,
- Fast Puiseux algorithm to be finished and written properly,
- Still one missing point for the  $\mathcal{O}(D^3)$  strategy,
- Efficient implementation of all these algorithms,
- Studying the next parts of the Abel map (optimizations, implementation. . . )